Perspective on the analysis and design of optimization algorithms

Adrien Taylor





AGIMUS Winter School — December 2023

Sierra team at Inria Paris



Main research themes: optimization, statistical learning, and interactions.

- ◊ Team leader: Francis Bach,
- ◊ Assistant: Marina Kovacic,
- ◊ 5 permanent researchers,
- ◊ 4 postdocs,
- ◊ 15 PhD students.

Perspective on the analysis and design of optimization algorithms

Adrien Taylor





AGIMUS Winter School — December 2023

 fixed-point perspective on optimization algorithms many references on the topic.

 ◊ fixed-point perspective on optimization algorithms many references on the topic.^{1,2,3,4}

¹Moreau (1962). "Fonctions convexes duales et points proximaux dans un espace hilbertien."

²Rockafellar (1976). "Augmented Lagrangians and applications of the proximal point algorithm in convex programming."

³Combettes, Pesquet (2011). "Proximal splitting methods in signal processing."

⁴Ryu, Boyd. (2016). "Primer on monotone operator methods."

fixed-point perspective on optimization algorithms
 many references on the topic.^{1,2,3,4}

◊ a few (standard) ideas for designing methods,

¹Moreau (1962). "Fonctions convexes duales et points proximaux dans un espace hilbertien."

²Rockafellar (1976). "Augmented Lagrangians and applications of the proximal point algorithm in convex programming."

³Combettes, Pesquet (2011). "Proximal splitting methods in signal processing."

⁴Ryu, Boyd. (2016). "Primer on monotone operator methods."

fixed-point perspective on optimization algorithms
 many references on the topic.^{1,2,3,4}

- ◊ a few (standard) ideas for designing methods,
- ◊ a constructive approach to algorithm analysis and design
 ... if time allows.

¹Moreau (1962). "Fonctions convexes duales et points proximaux dans un espace hilbertien."

²Rockafellar (1976). "Augmented Lagrangians and applications of the proximal point algorithm in convex programming."

³Combettes, Pesquet (2011). "Proximal splitting methods in signal processing."

⁴Ryu, Boyd. (2016). "Primer on monotone operator methods."



François Glineur



Julien Hendrickx



Etienne de Klerk



Ernest Ryu



Carolina Bergeling



Pontus Giselsson



Francis Bach



Jérôme Bolte



Yoel Drori



Alexandre d'Aspremont



Mathieu Barré



Radu Dragomir



Bryan Van Scoy



Laurent Lessard



Aymeric Dieuleveut



Céline Moucer



Baptiste Goujaud



Sebastian Banert



Manu Uphadyaya



Eduard Gorbunov



Gauthier Gidel



Antoine Bambade



Sarah Kazdadi



Justin Carpentier

$$f(x_{\star}) \triangleq \min_{x \in \mathcal{D}} f(x),$$

with f, \mathcal{D} convex.

$$f(x_{\star}) \triangleq \min_{x \in \mathcal{D}} f(x),$$

with f, \mathcal{D} convex.

Use **iterative algorithm** generating a sequence x_0, x_1, \ldots, x_N .

$$f(x_{\star}) \triangleq \min_{x \in \mathcal{D}} f(x),$$

with f, \mathcal{D} convex.

Use **iterative algorithm** generating a sequence x_0, x_1, \ldots, x_N .

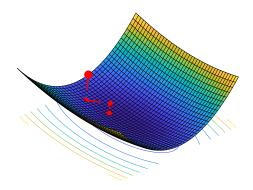
Example: gradient descent iterates $x_{k+1} = x_k - \gamma_k \nabla f(x_k)$.

$$f(x_{\star}) \triangleq \min_{x \in \mathcal{D}} f(x),$$

with f, \mathcal{D} convex.

Use **iterative algorithm** generating a sequence x_0, x_1, \ldots, x_N .

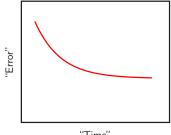
Example: gradient descent iterates $x_{k+1} = x_k - \gamma_k \nabla f(x_k)$.



Question: what a priori guarantees after N iterations?

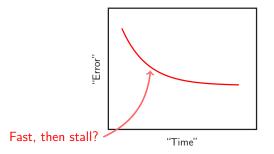
Question: what a priori guarantees after N iterations?

Question: what a priori guarantees after N iterations?

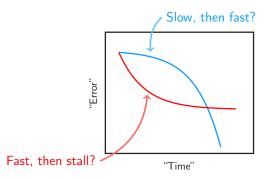


"Time"

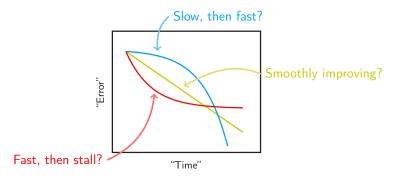
Question: what *a priori* guarantees after *N* iterations?



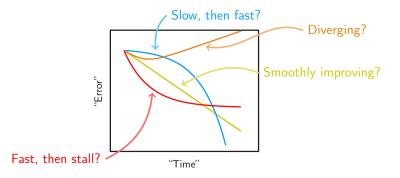
Question: what *a priori* guarantees after *N* iterations?



Question: what *a priori* guarantees after *N* iterations?



Question: what *a priori* guarantees after *N* iterations?



Classical approaches:

◊ "convergence" analysis (no speed),

- ◊ "convergence" analysis (no speed),
- ◇ "asymptotical" (local) analyses,

- ◊ "convergence" analysis (no speed),
- ◇ "asymptotical" (local) analyses,
- ◊ worst-case (global) analyses,

- ◊ "convergence" analysis (no speed),
- ◊ "asymptotical" (local) analyses,
- ◊ worst-case (global) analyses,
- ◊ average-case analyses,

- ◊ "convergence" analysis (no speed),
- ◊ "asymptotical" (local) analyses,
- ◊ worst-case (global) analyses,
- ◊ average-case analyses,
- ◊ high-probability analyses,

- ◊ "convergence" analysis (no speed),
- ◊ "asymptotical" (local) analyses,
- ◊ worst-case (global) analyses,
- ◊ average-case analyses,
- ◊ high-probability analyses,
- $\diamond\,$ smoothed analyses.

A few elements of convex analysis

We consider **extended-valued functions** $f : \mathbb{R}^n \to \mathbb{R} \cup \{+\infty\}$.

We consider **extended-valued functions** $f : \mathbb{R}^n \to \mathbb{R} \cup \{+\infty\}$.

Example: indicator function of interval [a, b]

$$i_{[a,b]}(x) = \left\{ egin{array}{cc} 0 & ext{if } x \in [a,b] \ +\infty & ext{otherwise.} \end{array}
ight.$$



We consider **extended-valued functions** $f : \mathbb{R}^n \to \mathbb{R} \cup \{+\infty\}$.

Example: indicator function of interval [a, b]

$$i_{[a,b]}(x) = \left\{ egin{array}{cc} 0 & ext{if } x \in [a,b] \ +\infty & ext{otherwise.} \end{array}
ight.$$

Effective domain of $f : \mathbb{R} \to \mathbb{R} \cup \{+\infty\}$:

dom $f = \{x : f(x) < +\infty\}.$

ā

We consider **extended-valued functions** $f : \mathbb{R}^n \to \mathbb{R} \cup \{+\infty\}$.

Example: indicator function of interval [a, b]

$$i_{[a,b]}(x) = \begin{cases} 0 & \text{if } x \in [a,b] \\ +\infty & \text{otherwise.} \end{cases}$$

Effective domain of $f : \mathbb{R} \to \mathbb{R} \cup \{+\infty\}$:

dom $f = \{x : f(x) < +\infty\}.$

ā

Note: so f might encode implicit constraints.

Convex functions

Graph is below line connecting any pairs (x, f(x)) and (y, f(y)):

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y),$$

 $\forall \theta \in [0, 1].$

Convex functions

Graph is below line connecting any pairs (x, f(x)) and (y, f(y)):

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y),$$

 $\forall \theta \in [0, 1].$

Examples:

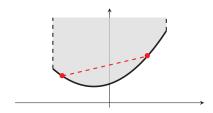
Convex functions

Graph is below line connecting any pairs (x, f(x)) and (y, f(y)):

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y),$$

 $\forall \theta \in [0,1].$

Examples:



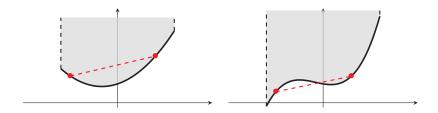
Convex functions

Graph is below line connecting any pairs (x, f(x)) and (y, f(y)):

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y),$$

 $\forall \theta \in [0,1].$

Examples:



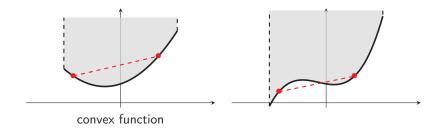
Convex functions

Graph is below line connecting any pairs (x, f(x)) and (y, f(y)):

$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y),$$

 $\forall \theta \in [0,1].$

Examples:



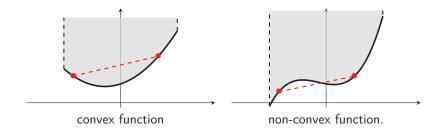
Convex functions

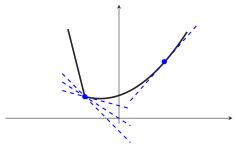
Graph is below line connecting any pairs (x, f(x)) and (y, f(y)):

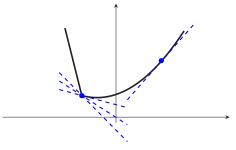
$$f(\theta x + (1 - \theta)y) \leq \theta f(x) + (1 - \theta)f(y),$$

 $\forall \theta \in [0,1].$

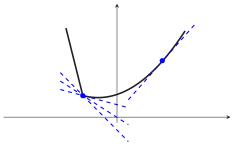
Examples:





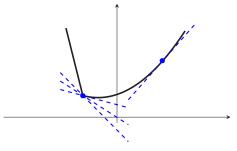


Notation:



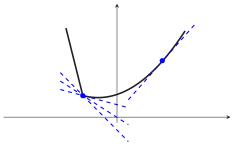
Notation:

♦ **subdifferential**: $\partial f : \mathbb{R}^n \to 2^{\mathbb{R}^n}$ (power set),



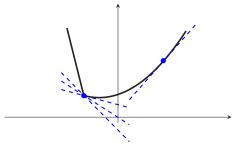
Notation:

- \diamond subdifferential: $\partial f : \mathbb{R}^n \to 2^{\mathbb{R}^n}$ (power set),
- ◊ subdifferential at x: $\partial f(x) = \{g : f(y) ≥ f(x) + g^T(y x) \forall y \in \mathbb{R}^n\},\$



Notation:

- ♦ **subdifferential**: $\partial f : \mathbb{R}^n \to 2^{\mathbb{R}^n}$ (power set),
- ◊ subdifferential at x: $\partial f(x) = \{g : f(y) ≥ f(x) + g^T(y x) \forall y \in \mathbb{R}^n\},\$
- ♦ any $g \in \partial f(x)$ is a **subgradient** of f at x.

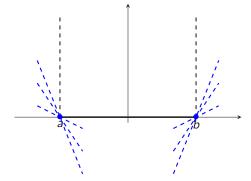


Notation:

- \diamond subdifferential: $\partial f : \mathbb{R}^n \to 2^{\mathbb{R}^n}$ (power set),
- ◊ subdifferential at x: $\partial f(x) = \{g : f(y) ≥ f(x) + g^T(y x) \forall y \in \mathbb{R}^n\},\$

♦ any $g \in \partial f(x)$ is a **subgradient** of f at x.

Essentially: (closed) convex functions have subgradients in relint(dom f).



Notation:

- \diamond subdifferential: $\partial f : \mathbb{R}^n \to 2^{\mathbb{R}^n}$ (power set),
- ◊ subdifferential at x: $\partial f(x) = \{g : f(y) ≥ f(x) + g^T(y x) \forall y \in \mathbb{R}^n\},\$

♦ any $g \in \partial f(x)$ is a **subgradient** of f at x.

Essentially: (closed) convex functions have subgradients in relint(dom f).

Optimality conditions

Convex optimization problem:

 $\min_{x\in\mathbb{R}^n}f(x)$

Optimality condition (Fermat's rule): x_* optimal iff $0 \in \partial f(x_*)$.

Optimality conditions

Convex optimization problem:

 $\min_{x\in\mathbb{R}^n}f(x)$

Optimality condition (Fermat's rule): x_* optimal iff $0 \in \partial f(x_*)$.

 \diamond Proof: x minimizes f if and only if

$$f(y) \geqslant f(x) = f(x) + 0^T (y - x)$$
 for all $y \in \mathbb{R}^n$.

Optimality conditions

Convex optimization problem:

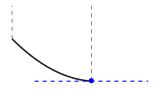
 $\min_{x\in\mathbb{R}^n}f(x)$

Optimality condition (Fermat's rule): x_* optimal iff $0 \in \partial f(x_*)$.

 \diamond Proof: x minimizes f if and only if

$$f(y) \ge f(x) = f(x) + 0^T (y - x)$$
 for all $y \in \mathbb{R}^n$.

♦ Example: several subgradients at solution, including 0:



Convex optimization problem (take II):

 $\min_{x\in\mathbb{R}^n}f(x)+h(x)$

with both f and h convex (closed, non-empty effective domains).

Convex optimization problem (take II):

 $\min_{x\in\mathbb{R}^n}f(x)+h(x)$

with both f and h convex (closed, non-empty effective domains).

 $\diamond~$ Under some conditions (e.g., "constraint qualifications"):

$$\partial (f+h)(x) = \partial f(x) + \partial h(x).$$

Convex optimization problem (take II):

 $\min_{x\in\mathbb{R}^n}f(x)+h(x)$

with both f and h convex (closed, non-empty effective domains).

 $\diamond~$ Under some conditions (e.g., "constraint qualifications"):

$$\partial (f+h)(x) = \partial f(x) + \partial h(x).$$

♦ We will search (Fermat's rule) for x_* :

$$0 \in \partial (f+h)(x_{\star}) = \partial f(x_{\star}) + \partial h(x_{\star}),$$

where equality follows from constraint qualifications (e.g., Slater's).

Convex optimization problem (take II):

 $\min_{x\in\mathbb{R}^n}f(x)+h(x)$

with both f and h convex (closed, non-empty effective domains).

 $\diamond~$ Under some conditions (e.g., "constraint qualifications"):

$$\partial (f+h)(x) = \partial f(x) + \partial h(x).$$

♦ We will search (Fermat's rule) for x_* :

$$0 \in \partial (f+h)(x_{\star}) = \partial f(x_{\star}) + \partial h(x_{\star}),$$

where equality follows from constraint qualifications (e.g., Slater's). This is strongly related to the usual KKT optimality conditions

Convex optimization problem (take II):

 $\min_{x\in\mathbb{R}^n}f(x)+h(x)$

with both f and h convex (closed, non-empty effective domains).

 $\diamond~$ Under some conditions (e.g., "constraint qualifications"):

$$\partial (f+h)(x) = \partial f(x) + \partial h(x).$$

♦ We will search (Fermat's rule) for x_* :

$$0 \in \partial (f+h)(x_{\star}) = \partial f(x_{\star}) + \partial h(x_{\star}),$$

where equality follows from constraint qualifications (e.g., Slater's). This is strongly related to the usual KKT optimality conditions

 \diamond in fact: it is exactly the same (alternate terminology)

Convex optimization problem (take II):

 $\min_{x\in\mathbb{R}^n}f(x)+h(x)$

with both f and h convex (closed, non-empty effective domains).

 $\diamond~$ Under some conditions (e.g., "constraint qualifications"):

$$\partial (f+h)(x) = \partial f(x) + \partial h(x).$$

♦ We will search (Fermat's rule) for x_* :

$$0 \in \partial(f+h)(x_{\star}) = \partial f(x_{\star}) + \partial h(x_{\star}),$$

where equality follows from constraint qualifications (e.g., Slater's). This is strongly related to the usual KKT optimality conditions

- ◊ in fact: it is exactly the same (alternate terminology)
- $\diamond~$ allows for simple manipulations.

A few approaches to constrained (convex) optimization

No free lunch: no "general" method works well for all problems.

No free lunch: no "general" method works well for all problems.

In numerical optimization, key is usually:

No free lunch: no "general" method works well for all problems.

In numerical optimization, key is usually:

◊ clear identification of my specifications (timings, accuracy, etc.)?

No free lunch: no "general" method works well for all problems.

In numerical optimization, key is usually:

- ◊ clear identification of my specifications (timings, accuracy, etc.)?
- ◊ divide problem into "cheap"/"simple" pieces...

No free lunch: no "general" method works well for all problems.

In numerical optimization, key is usually:

- ◊ clear identification of my specifications (timings, accuracy, etc.)?
- ◊ divide problem into "cheap"/"simple" pieces...

Current library of numerical optimization algorithms:

No free lunch: no "general" method works well for all problems.

In numerical optimization, key is usually:

- ◊ clear identification of my specifications (timings, accuracy, etc.)?
- ◊ divide problem into "cheap"/"simple" pieces...

Current library of numerical optimization algorithms:

◊ clearly a jungle,

No free lunch: no "general" method works well for all problems.

In numerical optimization, key is usually:

- $\diamond~$ clear identification of my specifications (timings, accuracy, etc.)?
- ◊ divide problem into "cheap"/"simple" pieces...

Current library of numerical optimization algorithms:

- ◊ clearly a jungle,
- ◊ still, a few key "algorithmic templates" emerged.

A few examples for the problem:

$$\min_{x\in\mathbb{R}^n} \left\{ f(x): x\in\mathcal{D} \right\}$$

with f a (closed, proper) convex function and \mathcal{D} a convex set.

A few examples for the problem:

$$\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}$$

with f a (closed, proper) convex function and \mathcal{D} a convex set.

A few examples for the problem:

$$\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}$$

with f a (closed, proper) convex function and D a convex set.

Typical splitting strategies:

 $\diamond~$ can I project on $\mathcal{D}?$

A few examples for the problem:

$$\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}$$

with f a (closed, proper) convex function and \mathcal{D} a convex set.

- $\diamond~$ can I project on $\mathcal{D}?$
- $\diamond~$ can I perform linear optimization on $\mathcal{D}?$

A few examples for the problem:

$$\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}$$

with f a (closed, proper) convex function and \mathcal{D} a convex set.

- $\diamond~$ can I project on $\mathcal{D}?$
- $\diamond~$ can I perform linear optimization on $\mathcal{D}?$
- ◊ can I compute gradients of f? Hessian?

A few examples for the problem:

$$\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}$$

with f a (closed, proper) convex function and \mathcal{D} a convex set.

- $\diamond~$ can I project on $\mathcal{D}?$
- $\diamond~$ can I perform linear optimization on $\mathcal{D}?$
- \diamond can I compute gradients of f? Hessian?
- ◊ stochastic approximations to f?

A few examples for the problem:

$$\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}$$

with f a (closed, proper) convex function and \mathcal{D} a convex set.

- $\diamond~$ can I project on $\mathcal{D}?$
- $\diamond~$ can I perform linear optimization on $\mathcal{D}?$
- \diamond can I compute gradients of f? Hessian?
- ◊ stochastic approximations to f?
- ◊ coordinate-wise optimization is easy?

A few examples for the problem:

$$\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}$$

with f a (closed, proper) convex function and \mathcal{D} a convex set.

Typical splitting strategies:

- $\diamond~$ can I project on $\mathcal{D}?$
- \diamond can I perform linear optimization on \mathcal{D} ?
- ◊ can I compute gradients of f? Hessian?
- ◊ stochastic approximations to f?
- ◊ coordinate-wise optimization is easy?

(choices also depends on the targets, e.g., accuracy).

Example: Projected gradient descent (I)

Differentiable (closed, proper) convex function f and convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}.$

Differentiable (closed, proper) convex function f and convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}.$

Assumptions:

Differentiable (closed, proper) convex function f and convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}.$

Assumptions:

 \diamond I can access gradients of f,

Differentiable (closed, proper) convex function f and convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}.$

Assumptions:

- \diamond I can access gradients of f,
- $\diamond \ \ \text{I can project on } \mathcal{D}.$

Differentiable (closed, proper) convex function f and convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}.$

Assumptions:

- \diamond I can access gradients of f,
- $\diamond \ \text{I can project on } \mathcal{D}.$

Iterate:
$$x_{k+1} = \operatorname*{argmin}_{x \in \mathcal{D}} \left\{ f(x_k) + \nabla f(x_k)^T (x - x_k) + \frac{2}{\gamma} \|x - x_k\|_2^2 \right\}.$$

Differentiable (closed, proper) convex function f and convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}\left\{f(x):\ x\in\mathcal{D}\right\}.$

Assumptions:

- \diamond I can access gradients of f,
- \diamond I can project on \mathcal{D} .

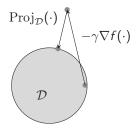
Iterate:
$$x_{k+1} = \operatorname*{argmin}_{x \in \mathcal{D}} \left\{ f(x_k) + \nabla f(x_k)^T (x - x_k) + \frac{2}{\gamma} \|x - x_k\|_2^2 \right\}.$$

Equivalently:

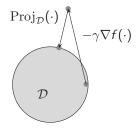
$$x_{k+1} = \operatorname*{argmin}_{x \in \mathcal{D}} \left\{ \|x - (x_k - \gamma \nabla f(x_k))\|_2^2 \right\}.$$

so $x_{k+1} = \operatorname{Proj}_{\mathcal{D}} (x_k - \gamma \nabla f(x_k)).$

 $x_{k+1} = \operatorname{Proj}_{\mathcal{D}} \left(x_k - \gamma \nabla f(x_k) \right)$



 $x_{k+1} = \operatorname{Proj}_{\mathcal{D}} \left(x_k - \gamma \nabla f(x_k) \right)$



Guarantees when f convex with L-Lipschitz gradient and $\gamma \in (0, 2/L)$. For instance, when $\gamma = 1/L$:

$$f(x_N) - f_\star \leqslant \frac{L \|x_0 - x_\star\|_2^2}{2N}$$

Smooth convex function f and (closed) convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}f(x)+i_{\mathcal{D}}(x),$

Smooth convex function f and (closed) convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}f(x)+i_{\mathcal{D}}(x),$

Smooth convex function f and (closed) convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}f(x)+i_{\mathcal{D}}(x),$

with (convex) indicator function of set \mathcal{D} : $i_{\mathcal{D}}(x) = \begin{cases} 0 & \text{if } x \in \mathcal{D} \\ +\infty & \text{otherwise.} \end{cases}$. Fixed-point viewpoint:

♦ we want to solve $0 \in \nabla f(x) + \partial i_D(x)$,

Smooth convex function f and (closed) convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}f(x)+i_{\mathcal{D}}(x),$

with (convex) indicator function of set \mathcal{D} : $i_{\mathcal{D}}(x) = \begin{cases} 0 & \text{if } x \in \mathcal{D} \\ +\infty & \text{otherwise.} \end{cases}$. Fixed-point viewpoint:

- \diamond we want to solve $0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x)$,
- ◊ base transformations:

 $0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x)$

Smooth convex function f and (closed) convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}f(x)+i_{\mathcal{D}}(x),$

- \diamond we want to solve $0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x)$,
- ◊ base transformations:

$$0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x) \Leftrightarrow 0 \in -\nabla f(x) - \partial i_{\mathcal{D}}(x)$$

Smooth convex function f and (closed) convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}f(x)+i_{\mathcal{D}}(x),$

- \diamond we want to solve $0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x)$,
- ◊ base transformations:

$$0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x) \Leftrightarrow 0 \in -\nabla f(x) - \partial i_{\mathcal{D}}(x)$$
$$\Leftrightarrow x \in x - \gamma \nabla f(x) - \gamma \partial i_{\mathcal{D}}(x)$$

Smooth convex function f and (closed) convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}f(x)+i_{\mathcal{D}}(x),$

- \diamond we want to solve $0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x)$,
- ◊ base transformations:

$$0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x) \Leftrightarrow 0 \in -\nabla f(x) - \partial i_{\mathcal{D}}(x)$$

$$\Leftrightarrow x \in x - \gamma \nabla f(x) - \gamma \partial i_{\mathcal{D}}(x)$$

$$\Leftrightarrow x + \gamma \partial i_{\mathcal{D}}(x) \ni x - \gamma \nabla f(x)$$

Smooth convex function f and (closed) convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}f(x)+i_{\mathcal{D}}(x),$

- \diamond we want to solve $0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x)$,
- ◊ base transformations:

$$0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x) \Leftrightarrow 0 \in -\nabla f(x) - \partial i_{\mathcal{D}}(x)$$

$$\Leftrightarrow x \in x - \gamma \nabla f(x) - \gamma \partial i_{\mathcal{D}}(x)$$

$$\Leftrightarrow x + \gamma \partial i_{\mathcal{D}}(x) \ni x - \gamma \nabla f(x)$$

$$\Leftrightarrow x = (I + \gamma \partial i_{\mathcal{D}})^{-1}(x - \gamma \nabla f(x))$$

Smooth convex function f and (closed) convex set \mathcal{D} :

 $\min_{x\in\mathbb{R}^n}f(x)+i_{\mathcal{D}}(x),$

- \diamond we want to solve $0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x)$,
- ◊ base transformations:

$$0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x) \Leftrightarrow 0 \in -\nabla f(x) - \partial i_{\mathcal{D}}(x)$$

$$\Leftrightarrow x \in x - \gamma \nabla f(x) - \gamma \partial i_{\mathcal{D}}(x)$$

$$\Leftrightarrow x + \gamma \partial i_{\mathcal{D}}(x) \ni x - \gamma \nabla f(x)$$

$$\Leftrightarrow x = (I + \gamma \partial i_{\mathcal{D}})^{-1}(x - \gamma \nabla f(x))$$

$$\Leftrightarrow x = \operatorname{Proj}_{\mathcal{D}}(x - \gamma \nabla f(x))$$

Differentiable (closed, proper) convex f and compact convex \mathcal{D} :

 $\min_{x\in\mathbb{R}^n} \left\{ f(x): x\in\mathcal{D} \right\}$

Differentiable (closed, proper) convex f and compact convex \mathcal{D} :

 $\min_{x\in\mathbb{R}^n} \left\{ f(x): x\in\mathcal{D} \right\}$

Assumptions:

Differentiable (closed, proper) convex f and compact convex \mathcal{D} :

 $\min_{x\in\mathbb{R}^n} \left\{ f(x): x\in\mathcal{D} \right\}$

Assumptions:

 \diamond I can access gradients of f,

Differentiable (closed, proper) convex f and compact convex \mathcal{D} :

 $\min_{x\in\mathbb{R}^n} \left\{ f(x): x\in\mathcal{D} \right\}$

Assumptions:

- \diamond I can access gradients of f,
- $\diamond\,$ I can perform linear optimization on $\mathcal{D}.$

Differentiable (closed, proper) convex f and compact convex \mathcal{D} :

 $\min_{x\in\mathbb{R}^n} \left\{ f(x): x\in\mathcal{D} \right\}$

Assumptions:

 \diamond I can access gradients of f,

 \diamond I can perform linear optimization on \mathcal{D} .

Iterate:

$$s_k = \operatorname*{argmin}_{s \in \mathcal{D}} \left\{ f(x_k) + \nabla f(x_k)^T (s - x_k) \right\}$$
$$x_k = (1 - \lambda_k) x_{k-1} + \lambda_k s_k.$$

Differentiable (closed, proper) convex f and compact convex \mathcal{D} :

 $\min_{x\in\mathbb{R}^n} \left\{ f(x): x\in\mathcal{D} \right\}$

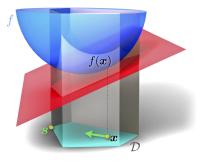
Assumptions:

- \diamond I can access gradients of f,
- \diamond I can perform linear optimization on \mathcal{D} .

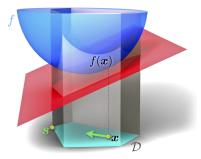
Iterate:

$$s_k = \operatorname*{argmin}_{s \in \mathcal{D}} \left\{ f(x_k) + \nabla f(x_k)^T (s - x_k) \right\}$$
$$x_k = (1 - \lambda_k) x_{k-1} + \lambda_k s_k.$$

Typical choice for $\lambda_k = \frac{2}{k+1}$.



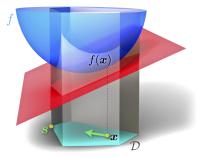
Picture from M. Jaggi (2013): "Revisiting Frank-Wolfe: Projection-Free Sparse Convex Optimization."



Picture from M. Jaggi (2013): "Revisiting Frank-Wolfe: Projection-Free Sparse Convex Optimization."

Guarantees when f convex with L-Lipschitz gradient and $Diam(\mathcal{D}) < \infty$:

$$f(x_N) - f_\star \leqslant \frac{L \mathrm{Diam}(\mathcal{D})^2}{N+2}$$



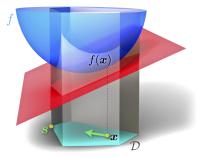
Picture from M. Jaggi (2013): "Revisiting Frank-Wolfe: Projection-Free Sparse Convex Optimization."

Guarantees when f convex with L-Lipschitz gradient and $Diam(\mathcal{D}) < \infty$:

$$f(x_N) - f_\star \leqslant \frac{L \mathrm{Diam}(\mathcal{D})^2}{N+2}$$

Similarly, can be seen as a fixed-point algorithm:

$$0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x)$$



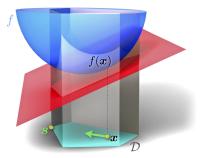
Picture from M. Jaggi (2013): "Revisiting Frank-Wolfe: Projection-Free Sparse Convex Optimization."

Guarantees when f convex with L-Lipschitz gradient and $Diam(\mathcal{D}) < \infty$:

$$f(x_N) - f_\star \leqslant \frac{L \mathrm{Diam}(\mathcal{D})^2}{N+2}$$

Similarly, can be seen as a fixed-point algorithm:

$$0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x) \Leftrightarrow x \in \partial i_{\mathcal{D}}^{-1}(-\nabla f(x))$$



Picture from M. Jaggi (2013): "Revisiting Frank-Wolfe: Projection-Free Sparse Convex Optimization."

Guarantees when f convex with L-Lipschitz gradient and $Diam(\mathcal{D}) < \infty$:

$$f(x_N) - f_{\star} \leqslant \frac{L \mathrm{Diam}(\mathcal{D})^2}{N+2}$$

Similarly, can be seen as a fixed-point algorithm:

$$0 \in \nabla f(x) + \partial i_{\mathcal{D}}(x) \Leftrightarrow x \in \partial i_{\mathcal{D}}^{-1}(-\nabla f(x))$$
$$\Leftrightarrow x \in (1 - \lambda)\partial i_{\mathcal{D}}^{-1}(-\nabla f(x)) + \lambda x.$$

- $\diamond~$ reformulate optimality conditions as fixed points,
 - using the operations that are "easy"
 - $-\,$ ex: can I invert a gradient? can I project? etc.

- reformulate optimality conditions as fixed points,
 - $-\,$ using the operations that are "easy"
 - $-\,$ ex: can I invert a gradient? can I project? etc.
- check/hope that the fixed-point iteration converges,

- reformulate optimality conditions as fixed points,
 - using the operations that are "easy"
 - ex: can I invert a gradient? can I project? etc.
- $\diamond~$ check/hope that the fixed-point iteration converges,
- $\diamond\,$ a few other standard schemes and many variations around them
 - Douglas-Rachford Splitting, ADMM, Chambolle-Pock algorithm, three-operator (Davis-Yin) splitting...

- reformulate optimality conditions as fixed points,
 - $-\,$ using the operations that are "easy"
 - ex: can I invert a gradient? can I project? etc.
- $\diamond~$ check/hope that the fixed-point iteration converges,
- $\diamond\,$ a few other standard schemes and many variations around them
 - Douglas-Rachford Splitting, ADMM, Chambolle-Pock algorithm, three-operator (Davis-Yin) splitting...
- $\diamond~$ most of those strategies apply to more general situations
 - $-\,$ that include primal-dual optimality conditions.

The proximal-point method

Continuous-time perspective

Minimization of a convex function f:

 $\min_{x\in\mathbb{R}^n}f(x)$

Continuous-time perspective

Minimization of a convex function f:

 $\min_{x\in\mathbb{R}^n}f(x)$

Continuous-time interpretation of the minimization procedure:

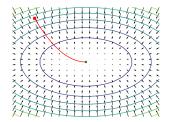
Continuous-time perspective

Minimization of a convex function f:

 $\min_{x\in\mathbb{R}^n}f(x)$

Continuous-time interpretation of the minimization procedure:

♦ "gradient flow": $\frac{d}{dt}x(t) = -\nabla f(x(t))$.



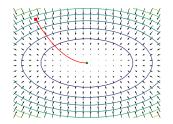
Continuous-time perspective

Minimization of a convex function f:

 $\min_{x\in\mathbb{R}^n}f(x)$

Continuous-time interpretation of the minimization procedure:

 \diamond "gradient flow": $\frac{d}{dt}x(t) = -\nabla f(x(t))$.



♦ Explicit Euler approx.: $x_{t+\Delta t} = x_t - \Delta t \nabla f(x_t)$ (gradient descent),

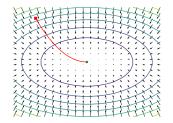
Continuous-time perspective

Minimization of a convex function f:

 $\min_{x\in\mathbb{R}^n}f(x)$

Continuous-time interpretation of the minimization procedure:

♦ "gradient flow": $\frac{d}{dt}x(t) = -\nabla f(x(t))$.



- ♦ Explicit Euler approx.: $x_{t+\Delta t} = x_t \Delta t \nabla f(x_t)$ (gradient descent),
- ♦ Implicit Euler approx.: $x_{t+\Delta t} = x_t \Delta t \nabla f(x_{t+\Delta t})$ (proximal point).

Minimize a (closed, proper) convex functions:

 $\min_{x\in\mathbb{R}^n}f(x)$

Minimize a (closed, proper) convex functions:

 $\min_{x\in\mathbb{R}^n}f(x)$

♦ Optimality condition: search x such that $0 \in \partial f(x)$.

Minimize a (closed, proper) convex functions:

 $\min_{x\in\mathbb{R}^n}f(x)$

- ♦ Optimality condition: search x such that $0 \in \partial f(x)$.
- ◊ A fixed-point reformulation:

$$x \in x - \gamma \partial f(x)$$

for some $\gamma > 0$.

Minimize a (closed, proper) convex functions:

 $\min_{x\in\mathbb{R}^n}f(x)$

- ♦ Optimality condition: search x such that $0 \in \partial f(x)$.
- ◊ A fixed-point reformulation:

$$x \in x - \gamma \partial f(x)$$

for some $\gamma > 0$.

Minimize a (closed, proper) convex functions:

 $\min_{x\in\mathbb{R}^n}f(x)$

- ♦ Optimality condition: search x such that $0 \in \partial f(x)$.
- ◊ A fixed-point reformulation:

$$x \in x - \gamma \partial f(x)$$

for some $\gamma > 0$.

♦ Proximal point: $x_{k+1} \in x_k - \gamma \partial f(x_{k+1})$.

Minimize a (closed, proper) convex functions:

 $\min_{x\in\mathbb{R}^n}f(x)$

- ♦ Optimality condition: search x such that $0 \in \partial f(x)$.
- ◊ A fixed-point reformulation:

$$x \in x - \gamma \partial f(x)$$

for some $\gamma > 0$.

♦ Proximal point: $x_{k+1} \in x_k - \gamma \partial f(x_{k+1})$. Equivalently:

$$x_{k+1} = \operatorname*{argmin}_{x \in \mathbb{R}^n} \left\{ f(x) + \frac{1}{2\gamma} \|x - x_k\|^2 \right\},$$

Minimize a (closed, proper) convex functions:

 $\min_{x\in\mathbb{R}^n}f(x)$

- ♦ Optimality condition: search x such that $0 \in \partial f(x)$.
- ◊ A fixed-point reformulation:

$$x \in x - \gamma \partial f(x)$$

for some $\gamma > 0$.

♦ Proximal point: $x_{k+1} \in x_k - \gamma \partial f(x_{k+1})$. Equivalently:

$$x_{k+1} = \operatorname*{argmin}_{x \in \mathbb{R}^n} \left\{ f(x) + \frac{1}{2\gamma} \|x - x_k\|^2 \right\},$$

 \diamond guaranteed:

$$f(x_k) - f(x_\star) \leq \frac{\|x_0 - x_\star\|^2}{4\sum_{i=0}^{k-1} \gamma_i}$$

Proximal-point operation:

 $\diamond~$ has a very long history

Proximal-point operation:

 \diamond has a very long history^{5,6,7}

⁵Moreau (1962). "Fonctions convexes duales et points proximaux dans un espace hilbertien."

⁶Minty (1962). "Monotone (nonlinear) operators in Hilbert space."

 7 Rockafellar (1976). "Augmented Lagrangians and applications of the proximal point algorithm in convex programming."

Proximal-point operation:

- \diamond has a very long history^{5,6,7}
- ◊ are at the center of many "first-order" algorithms,

⁵Moreau (1962). "Fonctions convexes duales et points proximaux dans un espace hilbertien."

⁶Minty (1962). "Monotone (nonlinear) operators in Hilbert space."

 7 Rockafellar (1976). "Augmented Lagrangians and applications of the proximal point algorithm in convex programming."

Proximal-point operation:

- ♦ has a very long history^{5,6,7}
- ♦ are at the center of many "first-order" algorithms,^{8,9,10}

⁵Moreau (1962). "Fonctions convexes duales et points proximaux dans un espace hilbertien."

⁶Minty (1962). "Monotone (nonlinear) operators in Hilbert space."

⁷Rockafellar (1976). "Augmented Lagrangians and applications of the proximal point algorithm in convex programming."

⁸Combettes, Pesquet (2011). "Proximal splitting methods in signal processing."

⁹Ryu, Boyd. (2016). "Primer on monotone operator methods."

 $^{10}{\rm Condat}$ (2022). "Proximal Splitting Algorithms for Convex Optimization: A Tour of Recent Advances, with New Twists."

Proximal-point operation:

- ♦ has a very long history^{5,6,7}
- ♦ are at the center of many "first-order" algorithms,^{8,9,10}
- $\diamond~$ Many examples of proximal operators with closed forms:

http://proximity-operator.net/,

⁵Moreau (1962). "Fonctions convexes duales et points proximaux dans un espace hilbertien."

⁶Minty (1962). "Monotone (nonlinear) operators in Hilbert space."

⁷Rockafellar (1976). "Augmented Lagrangians and applications of the proximal point algorithm in convex programming."

⁸Combettes, Pesquet (2011). "Proximal splitting methods in signal processing."

⁹Ryu, Boyd. (2016). "Primer on monotone operator methods."

 $^{10}{\rm Condat}$ (2022). "Proximal Splitting Algorithms for Convex Optimization: A Tour of Recent Advances, with New Twists."

Proximal-point operation:

- ♦ has a very long history^{5,6,7}
- ◊ are at the center of many "first-order" algorithms,^{8,9,10}
- Any examples of proximal operators with closed forms:

http://proximity-operator.net/,

many results on "approximated" proximal-point too.

⁵Moreau (1962). "Fonctions convexes duales et points proximaux dans un espace hilbertien."

⁶Minty (1962). "Monotone (nonlinear) operators in Hilbert space."

 7 Rockafellar (1976). "Augmented Lagrangians and applications of the proximal point algorithm in convex programming."

⁸Combettes, Pesquet (2011). "Proximal splitting methods in signal processing."

⁹Ryu, Boyd. (2016). "Primer on monotone operator methods."

 $^{10}{\rm Condat}$ (2022). "Proximal Splitting Algorithms for Convex Optimization: A Tour of Recent Advances, with New Twists."

Proximal-point for convex QPs (I)

Minimize a quadratic function (with $Q \succeq 0$) under linear constraints:

$$\min_{x} \left\{ \frac{1}{2} x^{T} Q x : A x \leqslant b \right\}.$$

with the proximal-point algorithm:

Proximal-point for convex QPs (I)

Minimize a quadratic function (with $Q \succeq 0$) under linear constraints:

$$\min_{x}\left\{\frac{1}{2}x^{T}Qx : Ax \leqslant b\right\}.$$

with the proximal-point algorithm:

$$x_{k+1} = \operatorname*{argmin}_{x \in \mathbb{R}^n} \left\{ f(x) + \frac{1}{2\gamma_k} \|x - x_k\|^2 \right\}.$$

Proximal-point for convex QPs (I)

Minimize a quadratic function (with $Q \geq 0$) under linear constraints:

$$\min_{x} \left\{ \frac{1}{2} x^{T} Q x : A x \leqslant b \right\}.$$

with the proximal-point algorithm:

$$x_{k+1} = \operatorname*{argmin}_{x \in \mathbb{R}^n} \left\{ f(x) + \frac{1}{2\gamma_k} \|x - x_k\|^2 \right\}.$$

Bad news:

$$x_{k+1} = \operatorname*{argmin}_{x \in \mathbb{R}^n} \left\{ \frac{1}{2} x^T Q x + \frac{1}{2\gamma_k} \| x - x_k \|^2 : A x \leqslant b \right\}$$

is just as hard as the original problem.

Minimize a quadratic function (with $Q \geq 0$) under linear constraints:

 $\min_{x} \left\{ \frac{1}{2} x^T Q x : A x \leqslant b \right\}.$

Minimize a quadratic function (with $Q \geq 0$) under linear constraints:

$$\min_{x}\left\{\frac{1}{2}x^{T}Qx : Ax \leqslant b\right\}.$$

Introduce $L(x, \lambda) = \frac{1}{2}x^TQx + q^Tx + \lambda^T(Ax - b)$

Minimize a quadratic function (with $Q \geq 0$) under linear constraints:

$$\min_{x}\left\{\frac{1}{2}x^{T}Qx : Ax \leqslant b\right\}.$$

Introduce $L(x, \lambda) = \frac{1}{2}x^T Q x + q^T x + \lambda^T (A x - b)$ and define dual $(\lambda \ge 0)$ $d(\lambda) = \min_{x} L(x, \lambda).$

Minimize a quadratic function (with $Q \geq 0$) under linear constraints:

$$\min_{x}\left\{\frac{1}{2}x^{T}Qx : Ax \leqslant b\right\}.$$

Introduce $L(x, \lambda) = \frac{1}{2}x^T Qx + q^T x + \lambda^T (Ax - b)$ and define dual $(\lambda \ge 0)$ $d(\lambda) = \min_x L(x, \lambda).$

 $d(\lambda)$ is concave so we can apply the proximal-point on the dual:

$$\lambda_{k+1} = \operatorname*{argmax}_{\lambda \geqslant 0} \left\{ d(\lambda) - rac{1}{2\gamma_k} \|\lambda - \lambda_k\|^2
ight\}.$$

Minimize a quadratic function (with $Q \succeq 0$) under linear constraints:

$$\min_{x}\left\{\frac{1}{2}x^{T}Qx : Ax \leqslant b\right\}.$$

Introduce $L(x, \lambda) = \frac{1}{2}x^T Qx + q^T x + \lambda^T (Ax - b)$ and define dual $(\lambda \ge 0)$ $d(\lambda) = \min_x L(x, \lambda).$

 $d(\lambda)$ is concave so we can apply the proximal-point on the dual:

$$\lambda_{k+1} = \operatorname*{argmax}_{\lambda \geqslant 0} \left\{ d(\lambda) - rac{1}{2\gamma_k} \|\lambda - \lambda_k\|^2
ight\}.$$

Explicitly max. over λ (not over x) yields the **method of mulipliers**:

$$x_{k+1} \in \operatorname*{argmin}_{x} L\left(x, [\lambda_k - \gamma_k(Ax - b)]_+\right)$$
$$\lambda_{k+1} = [\lambda_k - \gamma_k(Ax_{k+1} - b)]_+$$

(problem in x is piecewise quadratic, but convex and unconstrained).

Proximal-point also applies to the saddle-point formulation:

 $\max_{\lambda \geqslant 0} \min_{x} L(x,\lambda)$

which is convex-convave, yielding the proximal method of multipliers:

Proximal-point also applies to the saddle-point formulation:

 $\max_{\lambda \geqslant 0} \min_{x} L(x,\lambda)$

which is convex-convave, yielding the proximal method of multipliers:

$$(\lambda_{k+1}, x_{k+1}) = \operatorname*{argmax}_{\lambda \ge 0} \operatorname*{argmin}_{x} \left\{ L(x, \lambda) + \frac{1}{2\gamma_k} \|x - x_k\|^2 - \frac{1}{2\beta_k} \|\lambda - \lambda_k\|^2 \right\}.$$

Proximal-point also applies to the saddle-point formulation:

 $\max_{\lambda \geqslant 0} \min_{x} L(x,\lambda)$

which is convex-convave, yielding the proximal method of multipliers:

$$(\lambda_{k+1}, x_{k+1}) = \operatorname*{argmax}_{\lambda \ge 0} \operatorname*{argmin}_{x} \left\{ L(x, \lambda) + \frac{1}{2\gamma_k} \|x - x_k\|^2 - \frac{1}{2\beta_k} \|\lambda - \lambda_k\|^2 \right\}.$$

In practice, we need to:

◊ approximate

$$(\lambda_{k+1}, x_{k+1}) \approx_{\epsilon} \underset{\lambda \geqslant 0}{\operatorname{argmax}} \underset{x}{\operatorname{argmin}} \left\{ \mathcal{L}(x, \lambda) + \frac{1}{2\gamma_k} \|x - x_k\|^2 - \frac{1}{2\beta_k} \|\lambda - \lambda_k\|^2 \right\},$$

Proximal-point also applies to the saddle-point formulation:

 $\max_{\lambda \geqslant 0} \min_{x} L(x,\lambda)$

which is convex-convave, yielding the proximal method of multipliers:

$$(\lambda_{k+1}, x_{k+1}) = \operatorname*{argmax}_{\lambda \ge 0} \operatorname*{argmin}_{x} \left\{ L(x, \lambda) + \frac{1}{2\gamma_k} \|x - x_k\|^2 - \frac{1}{2\beta_k} \|\lambda - \lambda_k\|^2 \right\}.$$

In practice, we need to:

◊ approximate

$$(\lambda_{k+1}, x_{k+1}) \approx_{\epsilon} \underset{\lambda \geqslant 0}{\operatorname{argmax}} \underset{x}{\operatorname{argmin}} \left\{ \mathcal{L}(x, \lambda) + \frac{1}{2\gamma_k} \|x - x_k\|^2 - \frac{1}{2\beta_k} \|\lambda - \lambda_k\|^2 \right\},$$

 \diamond choose appropriate step size policies (β_k and γ_k) trading-off:

- large values: less iterations,
- small values: inner problem is simpler.



Antoine Bambade

Sarah Kazdadi



Fabian Schramm



Justin Carpentier

About

The Advanced Proximal Optimization Toolbox

cpp robotics optimization linear-programming proximal-algorithms quadratic-programming eigen3

- C Readme
- ধারু BSD-2-Clause license
- Cite this repository -
- -/- Activity
- ☆ 288 stars
- 14 watching
- 父 40 forks

Report repository

Releases 35



+ 34 releases

ProxSuite

License BSD 2-Clause docs online 💭 CI - Linux/OSX/Windows - Conda passing pypi package 0.3.6 Anaconda.org 0.3.6

ProxSuite is a collection of open-source, numerically robust, precise and efficient numerical solvers (e.g., LPs, QPs, etc.) rooted in revisited primal-dual proximal algorithms. Through **ProxSuite**, we aim to offer the community scalable optimizers that can deal with dense, sparse or matrix-free problems. While the first targeted application is Robotics, **ProxSuite** can be used in other contexts without limits.

ProxSuite is actively developped and supported by the Willow and Sierra research groups, joint research teams between Inria, École Normale Supérieure de Paris and Centre National de la Recherche Scientifique localized in France.

ProxSuite is already integrated into:

- · CVXPY modeling language for convex optimization problems,
- CasADI's symbolic framework for numerical optimization in general and optimal control. ProxQP is available in CasADI as plugin to solve quadratic programs,
- TSID: robotic software for efficient robot inverse dynamics with contacts and based on Pinocchio.

We are ready to integrate ProxSuite within other optimization ecosystems.

A priori working guarantees?

for numerical optimization algorithms

Classical approaches:

◊ "convergence" analysis (no speed),

- ◊ "convergence" analysis (no speed),
- ◊ "asymptotical" (local) analyses,

- ◊ "convergence" analysis (no speed),
- ◇ "asymptotical" (local) analyses,
- ◊ worst-case (global) analyses,

- ◊ "convergence" analysis (no speed),
- ◇ "asymptotical" (local) analyses,
- \diamond worst-case (global) analyses,
- ◊ average-case analyses,

- ◊ "convergence" analysis (no speed),
- ◇ "asymptotical" (local) analyses,
- \diamond worst-case (global) analyses,
- ◊ average-case analyses,
- ◊ high-probability analyses,

Convergence of an algorithm? (if it works)

Classical approaches:

- ◊ "convergence" analysis (no speed),
- ◊ "asymptotical" (local) analyses,
- \diamond worst-case (global) analyses,
- $\diamond~$ average-case analyses,
- ◊ high-probability analyses,
- $\diamond\,$ smoothed analyses.

Here: classical worst-case framework

Requires **assumptions** on f (formally: $f \in \mathcal{F}$ for a certain \mathcal{F}).

Requires **assumptions** on f (formally: $f \in \mathcal{F}$ for a certain \mathcal{F}).

How to ensure an algorithm to "always" work with prescribed guarantees?

Requires **assumptions** on f (formally: $f \in \mathcal{F}$ for a certain \mathcal{F}).

How to ensure an algorithm to "always" work with prescribed guarantees?

 \diamond "if it works on the **worst** $f \in \mathcal{F}$, it works on all $f \in \mathcal{F}$."

Requires **assumptions** on *f* (formally: $f \in \mathcal{F}$ for a certain \mathcal{F}).

How to ensure an algorithm to "always" work with prescribed guarantees?

 \diamond "if it works on the **worst** $f \in \mathcal{F}$, it works on all $f \in \mathcal{F}$."

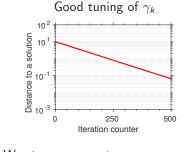
Philosophically, what we target:

Requires **assumptions** on f (formally: $f \in \mathcal{F}$ for a certain \mathcal{F}).

How to ensure an algorithm to "always" work with prescribed guarantees?

 \diamond "if it works on the **worst** $f \in \mathcal{F}$, it works on all $f \in \mathcal{F}$."

Philosophically, what we target:



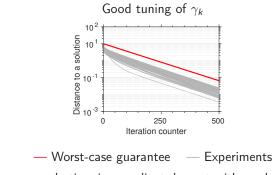
Worst-case guarantee

Requires **assumptions** on f (formally: $f \in \mathcal{F}$ for a certain \mathcal{F}).

How to ensure an algorithm to "always" work with prescribed guarantees?

 \diamond "if it works on the **worst** $f \in \mathcal{F}$, it works on all $f \in \mathcal{F}$."

Philosophically, what we target:



(Here: convex quadratic min., gradient descent with random init.)

Key points:

Key points:

 $\diamond~$ worst-case convergence \Rightarrow we can trust method as black-boxes,

Key points:

- $\diamond~$ worst-case convergence \Rightarrow we can **trust** method as black-boxes,
- $\diamond\,$ worst-case guarantees \Rightarrow guide for **method tuning**.

Key points:

- \diamond worst-case convergence \Rightarrow we can **trust** method as black-boxes,
- \diamond worst-case guarantees \Rightarrow guide for **method tuning**.

This is a whole field of study, with many results.^{11,12,13}

Key points:

- \diamond worst-case convergence \Rightarrow we can **trust** method as black-boxes,
- \diamond worst-case guarantees \Rightarrow guide for **method tuning**.

This is a whole field of study, with many results.^{11,12,13}

- A few limitations of the traditional viewpoint:
 - conservative nature,

Key points:

- $\diamond~$ worst-case convergence \Rightarrow we can trust method as black-boxes,
- \diamond worst-case guarantees \Rightarrow guide for **method tuning**.

This is a whole field of study, with many results.^{11,12,13}

- A few limitations of the traditional viewpoint:
 - conservative nature,
 - ◊ technical, error-prone,

Key points:

- $\diamond~$ worst-case convergence \Rightarrow we can trust method as black-boxes,
- \diamond worst-case guarantees \Rightarrow guide for **method tuning**.

This is a whole field of study, with many results.^{11,12,13}

- A few limitations of the traditional viewpoint:
 - conservative nature,
 - ◊ technical, error-prone,
 - lack global insights.

Global insight of worst-case analyses?

Find $x_{\star} \in \mathbb{R}^d$ such that

 $f(x_{\star}) = \min_{x \in \mathbb{R}^d} f(x).$

Find $x_{\star} \in \mathbb{R}^d$ such that

 $f(x_{\star}) = \min_{x \in \mathbb{R}^d} f(x).$

(Gradient method) We decide to use: $x_{k+1} = x_k - \gamma_k \nabla f(x_k)$

Find $x_{\star} \in \mathbb{R}^d$ such that

 $f(x_{\star}) = \min_{x \in \mathbb{R}^d} f(x).$

(Gradient method) We decide to use: $x_{k+1} = x_k - \gamma_k \nabla f(x_k)$

Question: what a priori guarantees after N iterations?

Find $x_{\star} \in \mathbb{R}^d$ such that

 $f(x_{\star}) = \min_{x \in \mathbb{R}^d} f(x).$

(Gradient method) We decide to use: $x_{k+1} = x_k - \gamma_k \nabla f(x_k)$

Question: what a priori guarantees after N iterations?

Examples: what about $f(x_N) - f(x_*)$, $\|\nabla f(x_N)\|$, $\|x_N - x_*\|$?

Find $x_{\star} \in \mathbb{R}^d$ such that

 $f(x_{\star}) = \min_{x \in \mathbb{R}^d} f(x).$

(Gradient method) We decide to use: $x_{k+1} = x_k - \gamma_k \nabla f(x_k)$

Question: what a priori guarantees after N iterations?

Examples: what about $f(x_N) - f(x_*)$, $\|\nabla f(x_N)\|$, $\|x_N - x_*\|$?

Alternatively: how fast does the ratio $\frac{\text{"error at iteration N"}}{\text{"initial error"}}$ decrease with N?

Find $x_{\star} \in \mathbb{R}^d$ such that

 $f(x_{\star}) = \min_{x \in \mathbb{R}^d} f(x).$

(Gradient method) We decide to use: $x_{k+1} = x_k - \gamma_k \nabla f(x_k)$

Question: what a priori guarantees after N iterations?

Examples: what about $f(x_N) - f(x_*)$, $\|\nabla f(x_N)\|$, $\|x_N - x_*\|$?

Alternatively: how fast does the ratio $\frac{\text{"error at iteration N"}}{\text{"initial error"}}$ decrease with N?

Such guarantees reachable only by assuming something on f.

Typical assumptions

Nontrivial guarantees only by assuming something on class of problems!

Typical assumptions

Nontrivial guarantees only by assuming something on class of problems!

Many standard (documented) classes of functions.

Typical assumptions

Nontrivial guarantees only by assuming something on class of problems!

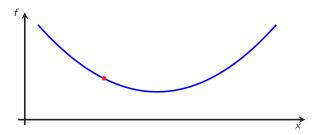
Many standard (documented) classes of functions.

Among many others: diff. function *f* is commonly assumed to be (for all $x, y \in \mathbb{R}^d$):

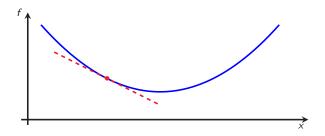
- $\diamond \quad L\text{-smooth: } f(x) \leqslant f(y) + \langle \nabla f(y), x y \rangle + \frac{L}{2} \|x y\|^2,$
- $\diamond \text{ convex: } f(x) \ge f(y) + \langle \nabla f(y), x y \rangle,$
- $\diamond \quad \mu\text{-strongly convex:} \ f(x) \geqslant f(y) + \langle \nabla f(y), x y \rangle + \frac{\mu}{2} \|x y\|^2.$

Consider a differentiable function $f : \mathbb{R}^d \to \mathbb{R}$, f is (μ -strongly) convex and L-smooth iff $\forall x, y \in \mathbb{R}^d$ we have:

Consider a differentiable function $f : \mathbb{R}^d \to \mathbb{R}$, f is (μ -strongly) convex and L-smooth iff $\forall x, y \in \mathbb{R}^d$ we have:

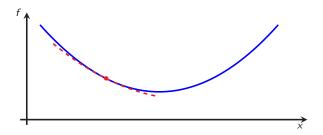


Consider a differentiable function $f : \mathbb{R}^d \to \mathbb{R}$, f is (μ -strongly) convex and L-smooth iff $\forall x, y \in \mathbb{R}^d$ we have:



(1) (Convexity) $f(x) \ge f(y) + \langle \nabla f(y), x - y \rangle$,

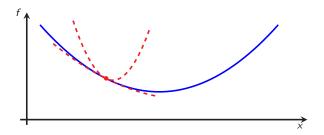
Consider a differentiable function $f : \mathbb{R}^d \to \mathbb{R}$, f is (μ -strongly) convex and L-smooth iff $\forall x, y \in \mathbb{R}^d$ we have:



(1) (Convexity) $f(x) \ge f(y) + \langle \nabla f(y), x - y \rangle$,

(1b) (μ -strong convexity) $f(x) \ge f(y) + \langle \nabla f(y), x - y \rangle + \frac{\mu}{2} ||x - y||^2$,

Consider a differentiable function $f : \mathbb{R}^d \to \mathbb{R}$, f is (μ -strongly) convex and L-smooth iff $\forall x, y \in \mathbb{R}^d$ we have:



(1) (Convexity)
$$f(x) \ge f(y) + \langle \nabla f(y), x - y \rangle$$
,

- (1b) (μ -strong convexity) $f(x) \ge f(y) + \langle \nabla f(y), x y \rangle + \frac{\mu}{2} ||x y||^2$,
 - (2) (L-smoothness) $f(x) \leq f(y) + \langle \nabla f(y), x y \rangle + \frac{l}{2} ||x y||^2$.

Toy example: What is the smallest τ such that:

$$||x_1 - x_\star||^2 \leq \tau ||x_0 - x_\star||^2$$

for all

Toy example: What is the smallest τ such that:

$$||x_1 - x_\star||^2 \leq \tau ||x_0 - x_\star||^2$$

for all

♦ *L*-smooth and μ -strongly convex function *f* (notation *f* ∈ $\mathcal{F}_{\mu,L}$),

Toy example: What is the smallest τ such that:

$$||x_1 - x_\star||^2 \leq \tau ||x_0 - x_\star||^2$$

for all

- ♦ *L*-smooth and μ -strongly convex function *f* (notation *f* ∈ $\mathcal{F}_{\mu,L}$),
- $\diamond x_0$, and x_1 generated by gradient step $x_1 = x_0 \gamma_0 \nabla f(x_0)$,

Toy example: What is the smallest τ such that:

$$||x_1 - x_\star||^2 \leq \tau ||x_0 - x_\star||^2$$

for all

- ♦ *L*-smooth and μ -strongly convex function *f* (notation *f* ∈ $\mathcal{F}_{\mu,L}$),
- $\diamond x_0$, and x_1 generated by gradient step $x_1 = x_0 \gamma_0 \nabla f(x_0)$,

$$x_{\star} = \operatorname*{argmin}_{x} f(x)?$$

Toy example: What is the smallest τ such that:

$$||x_1 - x_\star||^2 \leq \tau ||x_0 - x_\star||^2$$

for all

- ♦ *L*-smooth and μ -strongly convex function *f* (notation *f* ∈ $\mathcal{F}_{\mu,L}$),
- $\diamond x_0$, and x_1 generated by gradient step $x_1 = x_0 \gamma_0 \nabla f(x_0)$,

$$x_{\star} = \operatorname*{argmin}_{x} f(x)$$

Computing τ ?

Convergence rate of a gradient step

Toy example: What is the smallest τ such that:

$$||x_1 - x_\star||^2 \leq \tau ||x_0 - x_\star||^2$$

for all

♦ *L*-smooth and μ -strongly convex function *f* (notation *f* ∈ $\mathcal{F}_{\mu,L}$),

 $\diamond x_0$, and x_1 generated by gradient step $x_1 = x_0 - \gamma_0 \nabla f(x_0)$,

$$\diamond \ x_{\star} = \underset{x}{\operatorname{argmin}} \ f(x)$$

Computing τ ?

$$\tau = \max_{f, x_0, x_1, x_*} \frac{\|x_1 - x_*\|^2}{\|x_0 - x_*\|^2}$$

s.t. $f \in \mathcal{F}_{\mu, L}$ Functional class
 $x_1 = x_0 - \gamma_0 \nabla f(x_0)$ Algorithm
 $\nabla f(x_*) = 0$ Optimality of x_*

Convergence rate of a gradient step

Toy example: What is the smallest τ such that:

$$||x_1 - x_\star||^2 \leq \tau ||x_0 - x_\star||^2$$

for all

♦ *L*-smooth and μ -strongly convex function *f* (notation *f* ∈ $\mathcal{F}_{\mu,L}$),

 $\diamond x_0$, and x_1 generated by gradient step $x_1 = x_0 - \gamma_0 \nabla f(x_0)$,

$$\diamond \ x_{\star} = \underset{x}{\operatorname{argmin}} \ f(x)$$

Computing τ ?

$$\tau = \max_{f, x_0, x_1, x_*} \frac{\|x_1 - x_*\|^2}{\|x_0 - x_*\|^2}$$

s.t. $f \in \mathcal{F}_{\mu, L}$ Functional class
 $x_1 = x_0 - \gamma_0 \nabla f(x_0)$ Algorithm
 $\nabla f(x_*) = 0$ Optimality of x_*

<u>Variables</u>: f, x_0 , x_1 , x_* ; parameters: μ , L, γ_0 .

Convergence rate of a gradient step

Toy example: What is the smallest τ such that:

$$||x_1 - x_\star||^2 \leq \tau ||x_0 - x_\star||^2$$

for all

♦ *L*-smooth and μ -strongly convex function *f* (notation *f* ∈ $\mathcal{F}_{\mu,L}$),

♦ x_0 , and x_1 generated by gradient step $x_1 = x_0 - \gamma_0 \nabla f(x_0)$,

$$x_{\star} = \underset{x}{\operatorname{argmin}} f(x)$$

Computing τ ?

$$\tau = \max_{f, x_0, x_1, x_*} \frac{\|x_1 - x_*\|^2}{\|x_0 - x_*\|^2}$$

s.t. $f \in \mathcal{F}_{\mu, L}$ Functional class
 $x_1 = x_0 - \gamma_0 \nabla f(x_0)$ Algorithm
 $\nabla f(x_*) = 0$ Optimality of x_*

<u>Variables</u>: f, x_0 , x_1 , x_* ; parameters: μ , L, γ_0 . Problem can be reformulated as semidefinite program (SDP).

◊ Performance estimation problem:

$$\max_{\substack{f, x_0, x_1, x_* \\ subject \text{ to } }} \frac{\|x_1 - x_0\|^2}{\|x_0 - x_*\|^2}$$

subject to $f \text{ is } L\text{-smooth and } \mu\text{-strongly convex},$
 $x_1 = x_0 - \gamma_0 \nabla f(x_0)$
 $\nabla f(x_*) = 0.$

♦ Performance estimation problem:

$$\max_{\substack{f, x_0, x_1, x_* \\ \text{subject to}}} \frac{\|x_1 - x_0\|^2}{\|x_0 - x_*\|^2}$$

subject to f is L-smooth and μ -strongly convex,
 $x_1 = x_0 - \gamma_0 \nabla f(x_0)$
 $\nabla f(x_*) = 0.$

 \diamond Variables: f, x_0 , x_1 , x_{\star} .

♦ Performance estimation problem:

$$\begin{array}{ll} \max_{f, x_0, x_1, x_*} & \frac{\|x_1 - x_0\|^2}{\|x_0 - x_*\|^2} \\ \text{subject to} & f \text{ is } L\text{-smooth and } \mu\text{-strongly convex}, \\ & x_1 = x_0 - \gamma_0 \nabla f(x_0) \\ & \nabla f(x_*) = 0. \end{array}$$

- \diamond Variables: f, x_0 , x_1 , x_* .
- \diamond Sampled version: f is only used at x_0 and x_* (no need to sample other points)

◊ Performance estimation problem:

$$\begin{array}{ll} \max_{f, x_0, x_1, x_\star} & \frac{\|x_1 - x_0\|^2}{\|x_0 - x_\star\|^2} \\ \text{subject to} & f \text{ is } L\text{-smooth and } \mu\text{-strongly convex,} \\ & x_1 = x_0 - \gamma_0 \nabla f(x_0) \\ & \nabla f(x_\star) = 0. \end{array}$$

- \diamond Variables: f, x_0 , x_1 , x_{\star} .
- \diamond Sampled version: f is only used at x_0 and x_* (no need to sample other points)

$$\begin{array}{ll} \max_{\substack{x_0, x_1, x_* \\ g_0, g_* \\ f_0, f_* \end{array}} & \frac{\|x_1 - x_0\|^2}{\|x_0 - x_*\|^2} \\ \text{subject to} & \exists f \in \mathcal{F}_{\mu, L} \text{ such that } \begin{cases} f_i = f(x_i) & i = 0, * \\ g_i = \nabla f(x_i) & i = 0, * \\ x_1 = x_0 - \gamma_0 g_0 \\ g_* = 0. \end{cases}$$

◊ Performance estimation problem:

$$\begin{array}{ll} \max_{f,x_0,x_1,x_\star} & \frac{\|x_1 - x_0\|^2}{\|x_0 - x_\star\|^2} \\ \text{subject to} & f \text{ is L-smooth and μ-strongly convex,} \\ & x_1 = x_0 - \gamma_0 \nabla f(x_0) \\ & \nabla f(x_\star) = 0. \end{array}$$

- \diamond Variables: f, x_0 , x_1 , x_{\star} .
- \diamond Sampled version: f is only used at x_0 and x_* (no need to sample other points)

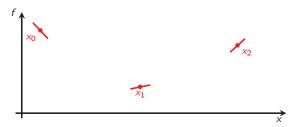
$$\max_{\substack{x_0, x_1, x_* \\ g_0, g_* \\ f_0, f_*}} \frac{\|x_1 - x_0\|^2}{\|x_0 - x_*\|^2}$$

subject to $\exists f \in \mathcal{F}_{\mu, L}$ such that
$$\begin{cases} f_i = f(x_i) & i = 0, \star \\ g_i = \nabla f(x_i) & i = 0, \star \\ x_1 = x_0 - \gamma_0 g_0 \\ g_\star = 0. \end{cases}$$

 \diamond Variables: x_0 , x_1 , x_{\star} , g_0 , g_{\star} , f_0 , f_{\star} .

Consider an index set S, and its associated values $\{(x_i, g_i, f_i)\}_{i \in S}$ with coordinates x_i , (sub)gradients g_i and function values f_i .

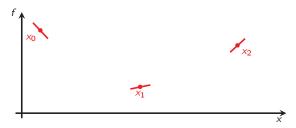
Consider an index set S, and its associated values $\{(x_i, g_i, f_i)\}_{i \in S}$ with coordinates x_i , (sub)gradients g_i and function values f_i .



? Possible to find $f \in \mathcal{F}_{\mu,L}$ such that

 $f(x_i) = f_i$, and $g_i = \nabla f(x_i)$, $\forall i \in S$.

Consider an index set S, and its associated values $\{(x_i, g_i, f_i)\}_{i \in S}$ with coordinates x_i , (sub)gradients g_i and function values f_i .



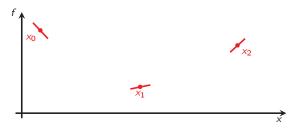
? Possible to find $f \in \mathcal{F}_{\mu,L}$ such that

 $f(x_i) = f_i$, and $g_i = \nabla f(x_i)$, $\forall i \in S$.

- Necessary and sufficient condition: $\forall i, j \in S$

$$f_i \ge f_j + \langle g_j, x_i - x_j \rangle + \frac{1}{2L} \|g_i - g_j\|^2 + \frac{\mu}{2(1-\mu/L)} \|x_i - x_j - \frac{1}{L}(g_i - g_j)\|^2.$$

Consider an index set S, and its associated values $\{(x_i, g_i, f_i)\}_{i \in S}$ with coordinates x_i , (sub)gradients g_i and function values f_i .



? Possible to find $f \in \mathcal{F}_{\mu,L}$ such that

$$f(x_i) = f_i$$
, and $g_i = \nabla f(x_i)$, $\forall i \in S$.

- Necessary and sufficient condition: $\forall i, j \in S$

$$f_i \ge f_j + \langle g_j, x_i - x_j \rangle + \frac{1}{2L} \|g_i - g_j\|^2 + \frac{\mu}{2(1-\mu/L)} \|x_i - x_j - \frac{1}{L}(g_i - g_j)\|^2$$

- Simpler example: pick $\mu = 0$ and $L = \infty$ (just convexity):

$$f_i \ge f_j + \langle g_j, x_i - x_j \rangle.$$

♦ Interpolation conditions allow removing red constraints

$$\max_{\substack{x_0, x_1, x_* \\ g_0, g_* \\ f_0, f_*}} \frac{\|x_1 - x_*\|^2}{\|x_0 - x_*\|^2}$$

subject to $\exists f \in \mathcal{F}_{\mu, L}$ such that
$$\begin{cases} f_i = f(x_i) & i = 0, * \\ g_i = \nabla f(x_i) & i = 0, * \\ x_1 = x_0 - \gamma_0 g_0 \\ g_* = 0, \end{cases}$$

♦ Interpolation conditions allow removing red constraints

$$\begin{array}{ll} \max_{\substack{x_{0}, x_{1}, x_{\star} \\ g_{0}, g_{\star}, f_{\star}, f_{\star} \\ f_{0}, f_{\star} \end{array}} & \frac{\|x_{1} - x_{\star}\|^{2}}{\|x_{0} - x_{\star}\|^{2}} \\ \text{subject to} & \exists f \in \mathcal{F}_{\mu, L} \text{ such that } \begin{cases} f_{i} = f(x_{i}) & i = 0, \star \\ g_{i} = \nabla f(x_{i}) & i = 0, \star \end{cases} \\ x_{1} = x_{0} - \gamma_{0} g_{0} \\ g_{\star} = 0, \end{cases}$$

◊ replacing them by

$$\begin{split} f_{\star} &\geq f_{0} + \langle g_{0}, x_{\star} - x_{0} \rangle + \frac{1}{2L} \|g_{\star} - g_{0}\|^{2} + \frac{\mu}{2(1-\mu/L)} \left\| x_{\star} - x_{0} - \frac{1}{L} (g_{\star} - g_{0}) \right\|^{2} \\ f_{0} &\geq f_{\star} + \langle g_{\star}, x_{0} - x_{\star} \rangle + \frac{1}{2L} \|g_{0} - g_{\star}\|^{2} + \frac{\mu}{2(1-\mu/L)} \left\| x_{0} - x_{\star} - \frac{1}{L} (g_{0} - g_{\star}) \right\|^{2}. \end{split}$$

◊ Interpolation conditions allow removing red constraints

$$\begin{array}{ll} \max_{\substack{x_{0}, x_{1}, x_{\star} \\ g_{0}, g_{\star}, f_{\star}, f_{\star} \\ f_{0}, f_{\star} \end{array}} & \frac{\|x_{1} - x_{\star}\|^{2}}{\|x_{0} - x_{\star}\|^{2}} \\ \text{subject to} & \exists f \in \mathcal{F}_{\mu, L} \text{ such that } \begin{cases} f_{i} = f(x_{i}) & i = 0, \star \\ g_{i} = \nabla f(x_{i}) & i = 0, \star \end{cases} \\ x_{1} = x_{0} - \gamma_{0} g_{0} \\ g_{\star} = 0, \end{cases}$$

◊ replacing them by

$$\begin{split} f_{\star} &\geq f_{0} + \langle g_{0}, x_{\star} - x_{0} \rangle + \frac{1}{2L} \|g_{\star} - g_{0}\|^{2} + \frac{\mu}{2(1-\mu/L)} \left\| x_{\star} - x_{0} - \frac{1}{L} (g_{\star} - g_{0}) \right\|^{2} \\ f_{0} &\geq f_{\star} + \langle g_{\star}, x_{0} - x_{\star} \rangle + \frac{1}{2L} \|g_{0} - g_{\star}\|^{2} + \frac{\mu}{2(1-\mu/L)} \left\| x_{0} - x_{\star} - \frac{1}{L} (g_{0} - g_{\star}) \right\|^{2}. \end{split}$$

♦ Same optimal value (no relaxation); but still non-convex quadratic problem.

 \diamond Using the new variables $G \succeq 0$ and F

$$G = \begin{bmatrix} \|x_0 - x_\star\|^2 & \langle g_0, x_0 - x_\star \rangle \\ \langle g_0, x_0 - x_\star \rangle & \|g_0\|^2 \end{bmatrix}, \quad F = f_0 - f_\star,$$

 \diamond Using the new variables $G \succeq 0$ and F

$$G = \begin{bmatrix} \|x_0 - x_\star\|^2 & \langle g_0, x_0 - x_\star \rangle \\ \langle g_0, x_0 - x_\star \rangle & \|g_0\|^2 \end{bmatrix}, \quad F = f_0 - f_\star,$$

 $\diamond~$ previous problem can be reformulated as a 2 \times 2 SDP

$$\begin{array}{ll} \max_{G,\,F} & G_{1,1} + \gamma_0^2 G_{2,2} - 2\gamma_0 G_{1,2} \\ \text{subject to} & F + \frac{L\mu}{2(L-\mu)} G_{1,1} + \frac{1}{2(L-\mu)} G_{2,2} - \frac{L}{L-\mu} G_{1,2} \leqslant 0 \\ & -F + \frac{L\mu}{2(L-\mu)} G_{1,1} + \frac{1}{2(L-\mu)} G_{2,2} - \frac{\mu}{L-\mu} G_{1,2} \leqslant 0 \\ & G_{1,1} = 1 \\ & G \succcurlyeq 0 \end{array}$$

 \diamond Using the new variables $G \succeq 0$ and F

$$G = \begin{bmatrix} \|x_0 - x_\star\|^2 & \langle g_0, x_0 - x_\star \rangle \\ \langle g_0, x_0 - x_\star \rangle & \|g_0\|^2 \end{bmatrix}, \quad F = f_0 - f_\star,$$

 $\diamond~$ previous problem can be reformulated as a 2 \times 2 SDP

$$\begin{array}{ll} \max_{G,\,F} & G_{1,1} + \gamma_0^2 \, G_{2,2} - 2\gamma_0 \, G_{1,2} \\ \text{subject to} & F + \frac{L\mu}{2(L-\mu)} \, G_{1,1} + \frac{1}{2(L-\mu)} \, G_{2,2} - \frac{L}{L-\mu} \, G_{1,2} \leqslant 0 \\ & -F + \frac{L\mu}{2(L-\mu)} \, G_{1,1} + \frac{1}{2(L-\mu)} \, G_{2,2} - \frac{\mu}{L-\mu} \, G_{1,2} \leqslant 0 \\ & G_{1,1} = 1 \\ & G \succcurlyeq 0 \end{array}$$

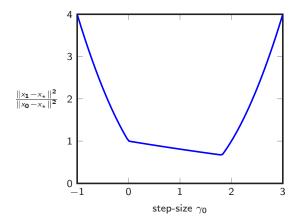
(using an an homogeneity argument and substituting x_1 and g_*).

Solving the SDP...

Fix L = 1, $\mu = .1$ and solve the SDP for a few values of γ_0 .

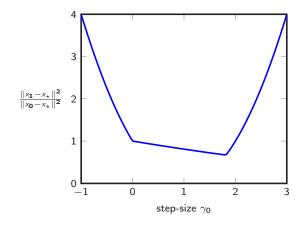
Solving the SDP...

Fix L = 1, $\mu = .1$ and solve the SDP for a few values of γ_0 .



Solving the SDP...

Fix L = 1, $\mu = .1$ and solve the SDP for a few values of γ_0 .



 $\diamond \ \ \, \text{Observation: numerics match max}\{(1-\gamma_0 L)^2,(1-\gamma_0 \mu)^2\}.$

 $\diamond~$ We can compute for the smallest $\tau(\gamma_{0})$ such that

$$||x_1 - x_*||^2 \leq \tau(\gamma_0) ||x_0 - x_*||^2$$

is satisfied for all $x_0 \in \mathbb{R}^d$, $d \in \mathbb{N}$, $f \in \mathcal{F}_{\mu,L}$, and $x_1 = x_0 - \gamma_0 \nabla f(x_0)$.

 $\diamond~$ We can compute for the smallest $\tau(\gamma_0)$ such that

$$||x_1 - x_*||^2 \leq \tau(\gamma_0) ||x_0 - x_*||^2$$

is satisfied for all $x_0 \in \mathbb{R}^d$, $d \in \mathbb{N}$, $f \in \mathcal{F}_{\mu,L}$, and $x_1 = x_0 - \gamma_0 \nabla f(x_0)$.

 \diamond Feasible points to the previous SDP correspond to lower bounds on $\tau(\gamma_0)$.

 $\diamond~$ We can compute for the smallest $\tau(\gamma_0)$ such that

$$||x_1 - x_\star||^2 \leq \tau(\gamma_0) ||x_0 - x_\star||^2$$

is satisfied for all $x_0 \in \mathbb{R}^d$, $d \in \mathbb{N}$, $f \in \mathcal{F}_{\mu,L}$, and $x_1 = x_0 - \gamma_0 \nabla f(x_0)$.

- \diamond Feasible points to the previous SDP correspond to lower bounds on $\tau(\gamma_0)$.
- ♦ Feasible points to dual SDP correspond to upper bounds on $\tau(\gamma_0)$.

 $\diamond~$ We can compute for the smallest $\tau(\gamma_0)$ such that

$$||x_1 - x_{\star}||^2 \leq \tau(\gamma_0) ||x_0 - x_{\star}||^2$$

is satisfied for all $x_0 \in \mathbb{R}^d$, $d \in \mathbb{N}$, $f \in \mathcal{F}_{\mu,L}$, and $x_1 = x_0 - \gamma_0 \nabla f(x_0)$.

- \diamond Feasible points to the previous SDP correspond to lower bounds on $\tau(\gamma_0)$.
- ♦ Feasible points to dual SDP correspond to upper bounds on $\tau(\gamma_0)$.
- ◊ Want to know more?
 - https://francisbach.com/computer-aided-analyses/
 - Toolboxes (next slides).

Software



François



Julien



Céline



Baptiste



Aymeric

- ◊ Performance Estimation Toolbox (PESTO) in Matlab, 2017.
- ◊ Performance Estimation in Python (PEPit), 2022.

Minimize *L*-smooth convex function f(x):

 $\min_{x\in\mathbb{R}^d}f(x).$

Minimize *L*-smooth convex function f(x):

$$\min_{x\in\mathbb{R}^d}f(x).$$

Accelerated Gradient Method Input: f L-smooth and convex, $x_0 = y_0 \in \mathbb{R}^d$. For i = 0 : N - 1 $x_{i+1} = y_i - \frac{1}{L} \nabla f(y_i)$ $y_{i+1} = x_{i+1} + \frac{i-1}{i+2} (x_{i+1} - x_i)$

Minimize *L*-smooth convex function f(x):

$$\min_{x\in\mathbb{R}^d}f(x).$$

Accelerated Gradient Method Input: f L-smooth and convex, $x_0 = y_0 \in \mathbb{R}^d$. For i = 0 : N - 1 $x_{i+1} = y_i - \frac{1}{L} \nabla f(y_i)$ $y_{i+1} = x_{i+1} + \frac{i-1}{i+2}(x_{i+1} - x_i)$

What if inexact gradient used instead? Relative inaccuracy model:

 $\|\widetilde{d}_f(y_i) - \nabla f(y_i)\| \leqslant \varepsilon \|\nabla f(y_i)\|.$

Minimize *L*-smooth convex function f(x):

$$\min_{x\in\mathbb{R}^d}f(x).$$

Accelerated Gradient Method Input: f L-smooth and convex, $x_0 = y_0 \in \mathbb{R}^d$. For i = 0 : N - 1 $x_{i+1} = y_i - \frac{1}{L} \nabla f(y_i)$ $y_{i+1} = x_{i+1} + \frac{i-1}{i+2} (x_{i+1} - x_i)$

What if inexact gradient used instead? Relative inaccuracy model:

 $\|\tilde{\mathsf{d}}_{\mathsf{f}}(\mathsf{y}_{\mathsf{i}}) - \nabla f(\mathsf{y}_{\mathsf{i}})\| \leq \varepsilon \|\nabla f(\mathsf{y}_{\mathsf{i}})\|.$

Minimize *L*-smooth convex function f(x):

$$\min_{x\in\mathbb{R}^d}f(x).$$

Accelerated Gradient Method Input: f L-smooth and convex, $x_0 = y_0 \in \mathbb{R}^d$. For i = 0 : N - 1 $x_{i+1} = y_i - \frac{1}{L} \tilde{d}_f(y_i)$ $y_{i+1} = x_{i+1} + \frac{i-1}{i+2}(x_{i+1} - x_i)$

What if inexact gradient used instead? Relative inaccuracy model:

 $\|\tilde{\mathsf{d}}_{\mathsf{f}}(\mathsf{y}_{\mathsf{i}}) - \nabla f(\mathsf{y}_{\mathsf{i}})\| \leq \varepsilon \|\nabla f(\mathsf{y}_{\mathsf{i}})\|.$

Minimize *L*-smooth convex function f(x):

$$\min_{x\in\mathbb{R}^d}f(x).$$

Accelerated Gradient Method Input: f L-smooth and convex, $x_0 = y_0 \in \mathbb{R}^d$. For i = 0 : N - 1 $x_{i+1} = y_i - \frac{1}{L} \widetilde{d}_f(y_i)$ $y_{i+1} = x_{i+1} + \frac{i-1}{i+2}(x_{i+1} - x_i)$

What if inexact gradient used instead? Relative inaccuracy model:

$$\|\widetilde{\mathsf{d}}_{\mathsf{f}}(\mathsf{y}_{\mathsf{i}}) - \nabla f(\mathsf{y}_{\mathsf{i}})\| \leq \varepsilon \|\nabla f(\mathsf{y}_{\mathsf{i}})\|.$$

What guarantees of type

$$\frac{f(x_N)-f_\star}{\|x_0-x_\star\|^2}\leqslant \tau(N,L)?$$

Next slide: compute $\tau(N, L)$ numerically using SDP.

```
% (0) Initialize an empty PEP
P = pep();
```

% (1) Set up the	objective function
param.mu = 0;	% strong convexity parameter
param.L = 1;	% Smoothness parameter

F=P.DeclareFunction('SmoothStronglyConvex', param); % F is the objective function

```
% (2) Set up the starting point and initial condition
x0 = P.StartingPoint(); % x0 is some starting point
[xs, fs] = F.OptimalPoint(); % xs is an optimal point, and fs=F(xs)
P.InitialCondition(\x0-xs\/2<= 1); % Add an initial condition |\x0-xs\/2<= 1</pre>
```

```
% (3) Algorithm
```

```
N = 7; % number of iterations
```

```
x = cell(N+1,1); % we store the iterates in a cell for convenience
x(1) = x0;
y = x0;
eps = .1;
for i = 1:N
d = inexactsubgradient(y, F, eps);
x(i+1) = y - 1/param.L * d;
y = x(i+1) + (i-1)/(i+2) * (x(i+1) - x(i));
end
```

```
% (4) Set up the performance measure
[g, f] = F.oracle(x{N+1}); % g=grad F(x), f=F(x)
P.PerformanceMetric(f - fs); % Worst-case evaluated as F(x)-F(xs)
```

```
% (5) Solve the PEP 
P.solve()
```

```
% (6) Evaluate the output
double(f - fs) % worst-case objective function accuracy
```

PESTO example: an inexact accelerated gradient method

% (0) Initialize an empty PEP P = pep();

% (1) Set up the objective function param.mu = 0; % strong convexity parameter param.L = 1; % Smoothness parameter

F=P.DeclareFunction('SmoothStronglyConvex',param); % F is the objective function

% (2) Set up the starting point and initial condition

```
x0 = P.StartingPoint(); % x0 is some starting point
fvs fcl = E Aotimal@oint(): % xs is an outlaal onint and fs=E(xs)
x{1} = x0;
y = x0;
eps = .1;
for i = 1:N
    d = inexactsubgradient(y, F, eps);
    x{i+1} = y - 1/param.L * d;
    y = x{i+1} + (i-1)/(i+2) * (x{i+1} - x{i});
end
    y = x(i+1) + (1-1)/(i+2) * (x[i+1] - x{i});
```

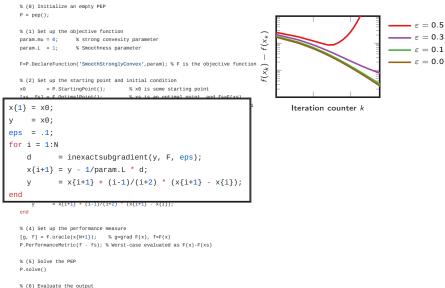
end

% (4) Set up the performance measure [g, f] = F.oracle(x{N+1}); % g=grad F(x), f=F(x) P.PerformanceMetric(f - fs); % Worst-case evaluated as F(x)-F(xs)

```
% (5) Solve the PEP 
P.solve()
```

% (6) Evaluate the output double(f - fs) % worst-case objective function accuracy

PESTO example: an inexact accelerated gradient method



double(f - fs) % worst-case objective function accuracy

Minimize sum of two convex (cpp) functions

 $\min_{x\in\mathbb{R}^d}f(x)+h(x).$

Douglas-Rachford Splitting Input: f, h convex (cpp) functions, $w_0 \in \mathbb{R}^d$. For i = 0 : N - 1 $x_{i+1} = \operatorname{argmin}_{x \in \mathbb{R}^d} \{\gamma h(x) + \frac{1}{2} ||x - w_i||^2\}$ $y_{i+1} = \operatorname{argmin}_{y \in \mathbb{R}^d} \{\gamma f(y) + \frac{1}{2} ||y - 2x_{i+1} + w_i||^2\}$ $w_{i+1} = w_i + \frac{1}{2}(y_{i+1} - x_{i+1}).$

Next slide: compute convergence rates when f is strongly convex and h is smooth.

```
% (O) Initialize an empty PEP
 P=pep();
 N = 1:
 % (1) Set up the class of monotone inclusions.
paramA.L = 1: paramA.mu = 0: % A is 1-Lipschitz and 0-strongly monotone
 paramB.mu = .1:
                               % B is .1-strongly monotone
 A = P.DeclareFunction('LipschitzStronglyMonotone',paramA);
 B = P.DeclareFunction('StronglyMonotone'.paramB);
 w = cell(N+1,1); wp = cell(N+1,1);
x = cell(N, 1); xp = cell(N, 1);
v = cell(N, 1); v_D = cell(N, 1);
% (2) Set up the starting points
w{1} = P.StartingPoint(): wp{1} = P.StartingPoint():
 P.InitialCondition((w{1}-wp{1})^2<=1);</pre>
% (3) Algorithm
 lambda = 1.3:
                  % step size (in the resolvents)
 theta = .9:
                    % overrelaxation
I for k = 1 : N
     x{k} = proximal step(w{k}.B.lambda):
            = proximal step(2*x{k}-w{k},A,lambda);
     v{k}
     w\{k+1\} = w\{k\} \cdot theta*(x\{k\} \cdot v\{k\}):
     xp{k} = proximal step(wp{k}.B.lambda);
     vp{k} = proximal_step(2*xp{k}-wp{k}.A.lambda);
     wp\{k+1\} = wp\{k\} \cdot theta*(xp\{k\} \cdot vp\{k\});
- end
% (4) Set up the performance measure: ||z0-z1||^2
 P.PerformanceMetric((w{k+1}-wp{k+1})^2):
 % (5) Solve the DED
 P.solve()
 % (6) Evaluate the output
 double((w{k+1}-wp{k+1})^2) % worst-case contraction factor
```

```
% (O) Initialize an empty PEP
P=pep():
N = 1:
% (1) Set up the class of monotone inclusions
paramA.L = 1: paramA.mu = 0: % A is 1-Lipschitz and 0-strongly monotone
paramB.mu = .1;
                              % B is .1-strongly monotone
A = P.DeclareFunction('LipschitzStronglyMonotone',paramA);
B = P.DeclareFunction('StronglyMonotone'.paramB);
w = cell(N+1,1); wp = cell(N+1,1);
x = cell(N, 1);
                 xp = cell(N, 1);
v = cell(N, 1); v_D = cell(N, 1);
% (2) Set up the starting points
w{1} = P.StartingPoint(): wp{1} = P.StartingPoint():
P.InitialCondition((w{1}-wp{1})^2<=1);</pre>
% (3) Algorithm
lambda = 1.3:
                    % step size (in the resolvents)
theta = .9:
                    % overrelaxation
            = proximal step(w{k},B,lambda);
 x{k
            = proximal step(2*x{k}-w{k},A,lambda);
 v{k}
            = w{k}-theta*(x{k}-v{k});
 w{k+1}
            = proximal step(wp(k), B.lambda):
    XD{k}
    vp{k}
            = proximal step(2*xp{k}-wp{k}.A.lambda):
    wp\{k+1\} = wp\{k\} \cdot theta*(xp\{k\} \cdot vp\{k\});
- end
% (4) Set up the performance measure: ||z0-z1||^2
P.PerformanceMetric((w{k+1}-wp{k+1})^2):
% (5) Solve the DED
P.solve()
% (6) Evaluate the output
double((w{k+1}-wp{k+1})^2) % worst-case contraction factor
```

% worst-case contraction factor

double((w{k+1}-wp{k+1})^2)

```
% (0) Initialize an empty PEP
P=pep():
N = 1:
% (1) Set up the class of monotone inclusions
paramA.L = 1: paramA.mu = 0: % A is 1-Lipschitz and 0-strongly monotone
paramB.mu = .1;
                               % B is .1-strongly monotone
                                                                  ٥
                                                                                                                  \mu = 0.1
A = P.DeclareFunction('LipschitzStronglyMonotone',paramA);
                                                                   Contraction rate
                                                                      0.8
B = P.DeclareFunction('StronglyMonotone'.paramB);
                                                                                                                  \mu = 0.5
                                                                      0.6
                                                                                                                  \mu = 1
w = cell(N+1.1):
                    wp = cell(N+1,1):
x = cell(N, 1):
                    xp = cell(N, 1);
                                                                                                                  \mu = 1.5
                                                                      0.4
v = cell(N, 1):
                    vp = cell(N, 1):
                                                                                                                  \mu = 2
                                                                      0.2
% (2) Set up the starting points
w{1} = P.StartingPoint(): wp{1}
                                     = P.StartingPoint():
P.InitialCondition((w{1}-wp{1})^2<=1);</pre>
                                                                                0.5
                                                                                               1.5
                                                                                                        2
                                                                                         1
                                                                              Lipschitz constant L
% (3) Algorithm
lambda = 1.3:
                    % step size (in the resolvents)
theta = .9:
                     % overrelaxation
 x{k}
            = proximal step(w{k},B,lambda);
            = proximal step(2*x{k}-w{k},A,lambda);
 v{k]
            = w{k}-theta*(x{k}-v{k});
 w{k+1}
            = proximal step(wp{k}.B.lambda):
     XD{k}
            = proximal step(2*xp{k}-wp{k},A,lambda);
     vp{k}
    wp{k+1}
              = wp\{k\} - theta*(xp\{k\} - vp\{k\}):
- end
% (4) Set up the performance measure: ||z0-z1||^2
P.PerformanceMetric((w{k+1}-wp{k+1})^2):
% (5) Solve the DED
P.solve()
% (6) Evaluate the output
```

Includes... but not limited to

 $\diamond~$ subgradient, gradient, heavy-ball, fast gradient, optimized gradient methods,

- $\diamond~$ subgradient, gradient, heavy-ball, fast gradient, optimized gradient methods,
- $\diamond~$ projected and proximal variants, accelerated/momentum versions,
- steepest descent, greedy/conjugate gradient methods,
- ◊ Frank-Wolfe/conditional gradients,

- $\diamond~$ subgradient, gradient, heavy-ball, fast gradient, optimized gradient methods,
- $\diamond~$ projected and proximal variants, accelerated/momentum versions,
- steepest descent, greedy/conjugate gradient methods,
- ◊ Frank-Wolfe/conditional gradients,
- ♦ Douglas-Rachford (ADMM), other operator splitting schemes,
- ◊ Krasnoselskii-Mann and Halpern fixed-point iterations,

- $\diamond~$ subgradient, gradient, heavy-ball, fast gradient, optimized gradient methods,
- $\diamond~$ projected and proximal variants, accelerated/momentum versions,
- steepest descent, greedy/conjugate gradient methods,
- ◊ Frank-Wolfe/conditional gradients,
- ♦ Douglas-Rachford (ADMM), other operator splitting schemes,
- ◊ Krasnoselskii-Mann and Halpern fixed-point iterations,
- ◊ inexact versions of all the above,
- $\diamond~$ stochastic versions: SGD, SAG, SAGA and variants.

Includes... but not limited to

- $\diamond~$ subgradient, gradient, heavy-ball, fast gradient, optimized gradient methods,
- $\diamond~$ projected and proximal variants, accelerated/momentum versions,
- steepest descent, greedy/conjugate gradient methods,
- ◊ Frank-Wolfe/conditional gradients,
- ♦ Douglas-Rachford (ADMM), other operator splitting schemes,
- ◊ Krasnoselskii-Mann and Halpern fixed-point iterations,
- ◊ inexact versions of all the above,
- $\diamond~$ stochastic versions: SGD, SAG, SAGA and variants.

Toolboxes contain most of the recent PEP-related advances (including techniques by other groups) available. Clean updated references & examples in user manual.

Includes... but not limited to

- $\diamond~$ subgradient, gradient, heavy-ball, fast gradient, optimized gradient methods,
- o projected and proximal variants, accelerated/momentum versions,
- steepest descent, greedy/conjugate gradient methods,
- ◊ Frank-Wolfe/conditional gradients,
- ♦ Douglas-Rachford (ADMM), other operator splitting schemes,
- ◊ Krasnoselskii-Mann and Halpern fixed-point iterations,
- ◊ inexact versions of all the above,
- $\diamond~$ stochastic versions: SGD, SAG, SAGA and variants.

Toolboxes contain most of the recent PEP-related advances (including techniques by other groups) available. Clean updated references & examples in user manual.

First ideas in this line of research coined by Drori and Teboulle (2014).

PEPit: Performance Estimation in Python

C Tests passing Codecov 90% docs passing pypi package 0.2.1 downloads 12k license MIT

This open source Python ilibrary provides a generic way to use PEP framework in Python. Performance estimation problems were introduced in 2014 by Yoel Drori and Mar Teboulle, see [11]. PEPt is mainly based on the formalism and developments from [2, 3] by a subset of the authors of this toolbox. A friendly informal introduction to this formalism is available in this blog post and a corresponding Matlab library is presented in [4] (PESTO).

Website and documentation of PEPit: https://pepit.readthedocs.io/

Source Code (MIT): https://github.com/PerformanceEstimation/PEPit

Using and citing the toolbox

This code comes jointly with the following reference :

B. Goujaud, C. Moucer, F. Glineur, J. Hendrickx, A. Taylor, A. Dieuleveut (2022). "PEPit: computer-assisted worst-case analyses of first-order optimization methods in Python."

When using the toolbox in a project, please refer to this note via this Bibtex entry:

@article{pepit2022,

tille=((PEP1t): computer-assisted worst-case analyses of first-order optimization methods in (P)yt author=(Goulant author=(

3

About

PEPit is a package enabling computerassisted worst-case analyses of firstorder optimization methods.

pepit.readthedocs.io/en/latest/

python optimization semidefinite-programming worst-case-analyses first-order-methods

performance-estimation-problems

Ш	Readme
ъ	MIT license
-^-	Activity
☆	68 stars
۲	4 watching
ę	6 forks
Report repository	

Releases 5



+ 4 releases

Important inspiration & reference:

Orori, and Teboulle ('14). "Performance of first-order methods for smooth convex minimization: a novel approach." Important inspiration & reference:

 Drori, and Teboulle ('14). "Performance of first-order methods for smooth convex minimization: a novel approach."

Second part of the presentation:

- ◊ T., Hendrickx, Glineur ('17). "Smooth strongly convex interpolation and exact worst-case performance of first-order methods."
- T., Hendrickx, Glineur ('17). "Exact worst-case performance of first-order methods for composite convex optimization."
- ◊ T., Hendrickx, Glineur ('17). "Performance estimation toolbox (PESTO): Automated worst-case analysis of first-order optimization methods."
- Goujaud, Moucer, et al. ('22). "PEPit: computer-assisted worst-case analyses of first-order optimization methods in Python."

Important inspiration & reference:

◊ Drori, and Teboulle ('14). "Performance of first-order methods for smooth convex minimization: a novel approach."

Second part of the presentation:

- ◊ T., Hendrickx, Glineur ('17). "Smooth strongly convex interpolation and exact worst-case performance of first-order methods."
- T., Hendrickx, Glineur ('17). "Exact worst-case performance of first-order methods for composite convex optimization."
- ◊ T., Hendrickx, Glineur ('17). "Performance estimation toolbox (PESTO): Automated worst-case analysis of first-order optimization methods."
- Goujaud, Moucer, et al. ('22). "PEPit: computer-assisted worst-case analyses of first-order optimization methods in Python."

Designing algorithms with PEPs:

- ◊ Drori, T ('20). "Efficient first-order methods for convex minimization: a constructive approach."
- $\diamond~$ Drori, T ('22). "On the oracle complexity of smooth strongly convex minimization."
- ◊ T, Drori ('23). "An optimal gradient method for smooth strongly convex minimization."

Informal introduction: https://francisbach.com/computer-aided-analyses/.

Optimization algorithms: currently a wild jungle.

- $\diamond~$ still: certain guiding principles & main driving algorithms,
- $\diamond~$ guarantees \Rightarrow trust and black box,
- $\diamond \ \ldots$ but somewhat behind schedule.

Optimization algorithms: currently a wild jungle.

- $\diamond~$ still: certain guiding principles & main driving algorithms,
- $\diamond~$ guarantees \Rightarrow trust and black box,
- $\diamond \ \ldots$ but somewhat behind schedule.

Optimization algorithms: currently a wild jungle.

- $\diamond~$ still: certain guiding principles & main driving algorithms,
- $\diamond~$ guarantees \Rightarrow trust and black box,
- $\diamond \ \ldots$ but somewhat behind schedule.

- ◊ numerically allows obtaining tight bounds (rigorous baselines),
 - fast prototyping
 - $-\,$ worth checking before trying to prove a method works.

Optimization algorithms: currently a wild jungle.

- $\diamond~$ still: certain guiding principles & main driving algorithms,
- $\diamond~$ guarantees \Rightarrow trust and black box,
- $\diamond \ \ldots$ but somewhat behind schedule.

- ◊ numerically allows obtaining tight bounds (rigorous baselines),
 - fast prototyping
 - worth checking before trying to prove a method works.
- ◊ algebraic insights into proofs: principled approach.

Optimization algorithms: currently a wild jungle.

- $\diamond~$ still: certain guiding principles & main driving algorithms,
- $\diamond~$ guarantees \Rightarrow trust and black box,
- $\diamond \ \ldots$ but somewhat behind schedule.

- ◊ numerically allows obtaining tight bounds (rigorous baselines),
 - fast prototyping
 - worth checking before trying to prove a method works.
- ◊ algebraic insights into proofs: principled approach.
- \diamond validation & benchmark tool for proofs (also for reviews s).

Thanks! Questions?

PerformanceEstimation/Performance-Estimation-Toolbox on Github

 $\operatorname{PerformanceEstimation}/\operatorname{PEPit}$ on Github